MEASURING UNCERTAINTY WITHOUT A NORM

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Abstract

Traub, Wasilkowski, and Woźniakowski have shown how uncertainty can be defined and analyzed without a norm or metric. Thier theory is based on two natural and non-restrictive axioms. We show that these axioms induce a family of pseudometrics, and that balls of radius ϵ are (roughly) the ϵ -approximations to the solution. In addition, we show that a family of pseudometrics is necessary, even for the problem of computing x such that $|f(x)| \leq \epsilon$, where f is a real function.

1. Introduction

In two recent monographs ([3],[4]), Traub and his colleagues have studied the optimal solution of problems which are solved approximately, that is, where there is uncertainty in the answer. In [4], uncertainty was measured by a norm. For some problems, this is not an appropriate or natural assumption. Therefore, in [3] it is shown how uncertainty can be introduced via two natural and non-restrictive axioms.

In a private communication, Traub asked about the strength of these axioms. That is, do the axioms generate any interesting structures? In Section 2 of this paper, we show that these axioms induce a family of pseudometrics. Moreover, we show that the balls of radius ε generated by this family of pseudometrics are (roughly speaking) the ε -approximations to the solution.

Is a <u>family</u> of pseudometrics necessary? We give an affirmative answer in Section 3, using the problem of computing x such that $|f(x)| \le \epsilon$, where f is a real function.

2. Solution Operators Are Generated by Families of Pseudometrics

We first recall the definition of a solution operator from [3]. Let F and G be sets, and let 2^G denote the power set of G, i.e., the class of all subsets of G. Let \mathbb{R}^+ denote the nonnegative real numbers. If $S: F \times \mathbb{R}^+ \to 2^G$ is an operator such that

$$(2.1) \qquad \forall f \in F, \qquad S(f,0) \neq \emptyset$$

and

$$(2.2) \quad \forall \ f \in F, \ \forall \ \epsilon_1, \epsilon_2 \in \mathbb{R}^+ \ \text{with} \ \epsilon_1 \leq \epsilon_2, \ S(f, \epsilon_1) \subseteq S(f, \epsilon_2),$$

then S is said to be a <u>solution</u> <u>operator</u>, and $S(f,\epsilon)$ is said to be the set of ϵ -approximations to the (exact) solution S(f,0).

Note that $S(f,\varepsilon)$ is a set. This formulation allows the exact solution S(f,0) to be a set, i.e., a problem may have multiple solutions. In addition, $S(f,\varepsilon)$ a set means that we are willing to accept any element of $S(f,\varepsilon)$ as an ε -approximation. These axioms are very natural: the first says that every problem has a solution, while the second says that increasing the uncertainty cannot decrease the family of ε -approximations.

In order to clarify these notions we give three examples.

Example 2.1. Let F be a set and let G be a normed linear space. Let $\overline{S}:F\to G$ be an operator. Define $S:F\times \mathbb{R}^+\to 2^G$ by

$$S(f,\epsilon) = \{g \in G : \|\overline{S}f - g\| \leq \epsilon\}.$$

Then S is a solution operator, and g \in G is an $_{\epsilon}$ -approximation

to $\overline{S}f$ precisely when $\|g - \overline{S}f\| \le \epsilon$. (This is the setting extensively studied in [4].)

Example 2.2. For a continuous function $f : \mathbb{R} \to \mathbb{R}$, let $Z(f) := \{x \in \mathbb{R} : f(x) = 0\}$

denote the zeroset of f. Now let

 $F = \{f : \mathbb{R} \to \mathbb{R} | f \text{ is continuous and } \mathbb{Z}(f) \neq \emptyset\},$

choose $G = \mathbb{R}$, and define $S : F \times \mathbb{R}^+ \rightarrow 2^G$ by

$$(2.3) S(f, s) := \{x \in \mathbb{R} : |f(x)| \leq \epsilon\}.$$

Then S is a solution operator, and $x \in S(f, \epsilon)$ precisely when $|f(x)| \leq \epsilon$, i.e., the residual of f at x is at most ϵ .

Before introducing the last example, we recall that a <u>pseudo-metric</u> d <u>on</u> G is a map d : $G \times G \rightarrow IR^+$ satisfying

$$d(g,g) = 0 Y g \in G,$$

$$(2.4) d(g_1,g_2) = d(g_2,g_1) Y g_1,g_2 \in G,$$

$$d(g_1,g_3) \leq d(g_1,g_2) + d(g_2,g_3) Y g_1,g_2,g_3 \in G.$$

(This terminology is standard in topology, see [2, pg. 198]. However, Collatz [1, pg. 21] refers to such a map as a "quasimetric," letting "pseudometric" refer to another concept entirely [1, pg. 51].) If d is a pseudometric on G, the set

$$B(g,d,\epsilon) := \{x \in G : d(x,g) \leq \epsilon\}$$

is called the d-ball of radius ϵ centered at g.

Example 2.3. Let F and G be sets. Let \emptyset be an F-indexed family of pseudometrics on G. Let $\overline{S}:F\to G$ be an operator. For $f\in F$, choose $d_f\in \emptyset$, and define

$$S(f, \varepsilon) := B(\overline{S}f, d_f, \varepsilon) \quad \forall \quad \varepsilon \geq 0.$$

Then it is easy to see that $S: F \times \mathbb{R}^+ \rightarrow 2^G$ is a solution operator.

We now show that, roughly speaking, Example 3 is the most general example of a solution operator.

Theorem 2.1. Let F,G be sets, and let $S: F \times \mathbb{R}^+ \to 2^G$ be a solution operator. Then for any $\epsilon_0 > 0$, there is an operator $\overline{S}: F \to G$ and a family $\mathfrak{D} = \{d_{\mathbf{f}}: \mathbf{f} \in F\}$ of pseudometrics on G such that

(2.5)
$$S(f, \varepsilon) \subseteq B(\overline{S}f, d_f, \varepsilon) \subseteq S(f, \varepsilon')$$

for any $f \in F$, any $\varepsilon \in [0, \varepsilon_0]$, and any $\varepsilon' \in (\varepsilon, \varepsilon_0]$. Proof: Let $f \in F$. For $g \in G$, let

$$D_{f}(g) := \left\{ \epsilon \geq 0 : g \in S(f, \epsilon) \right\},\$$

and now define $d_f : G \times G \rightarrow \mathbb{R}^+$ by

(2.6)
$$d_f(g_1,g_2) := \min\{\epsilon_0, |\inf D_f(g_2) - \inf D_f(g_1)|\},$$

where the inf of an empty set is defined to be ∞ , and ∞ - ∞ = 0. We first show that d_f is a pseudometric on G. Clearly, the

first two properties in (2.4) hold for d_f ; we need only check the third (the triangle inequality). Let $g_1, g_2, g_3 \in G$, and let

 $\delta_{i} = \inf D_{f}(g_{i})$. Then

$$|\delta_{3} - \delta_{1}| \leq |\delta_{2} - \delta_{1}| + |\delta_{3} - \delta_{2}|.$$

Arguing by cases if necessary, it is easy to see that (2.6) and (2.7) yield the triangle inequality.

We now define $\overline{S} : F \rightarrow G$ to be any map such that

$$\overline{S}f \in S(f,0)$$
 $\forall f \in F.$

This is possible because $S(f,0) \neq \emptyset$.

We now must prove (2.5). Let $f\in F$ and $\epsilon\in[0,\epsilon_0^{})\,.$ We first claim that

(2.8)
$$d_{f}(g,\overline{S}f) = \min\{\epsilon_{0}, \inf D_{f}(g)\}.$$

Indeed, since $\overline{S}f \in S(f,0)$, we have

inf
$$D_f(\overline{S}f) = 0$$
,

so that (2.8) follows from (2.6).

To see that $S(f,\epsilon)\subseteq B(\overline{S}f,d_f,\epsilon)$, let $g\in S(f,\epsilon)$. We then have

inf
$$D_f(g) \leq \epsilon$$
;

since ϵ $\langle \epsilon_0$, we use (2.8) to find

$$d_f(g,\overline{S}f) \leq \varepsilon$$
,

and so $g \in B(\overline{S}f,d_{f},\epsilon)$.

Now let $\varepsilon' \in (\varepsilon, \varepsilon_0]$. We wish to show that $B(\overline{S}f, d_f, \varepsilon) \subseteq S(f, \varepsilon')$. Let $g \in B(\overline{S}f, d_f, \varepsilon)$. Since $\varepsilon < \varepsilon_0$, we use (2.8) to find

(2.9)
$$\inf D_{f}(g) = d_{f}(g, \overline{S}f) \leq \varepsilon.$$

The first part of (2.9) and the definition of infimum yield

$$g \in S(f,d_f(g,\overline{S}f) + \delta)$$

for all $\delta>0$, no matter how small. Setting $\delta=\epsilon'-\epsilon>0$, we then have

$$g \in S(f, d_f(g, \overline{S}f) + \epsilon' - \epsilon)$$

$$\subseteq S(f, \epsilon + \epsilon' - \epsilon)$$

$$= S(f, \epsilon').$$

where the inclusion follows from the second part of (2.9) and the monotonicity condition (2.2).

Remark 2.1. We comment on the role played by ε_0 . It is possible to describe problems with $D_f(g)$ empty for some $f \in F$ and $g \in G$. (For example, take $\overline{S} : F \to G$ to be any operator where G is not a singleton, and define $S(f,\varepsilon) := \{\overline{S}f\} \ \forall \ f \in F$, $\varepsilon \geq 0$; then for any $f \in F$, $g \in G$ with $g \neq \overline{S}f$, and $\varepsilon \geq 0$, $g \notin S(f,\varepsilon)$, so that $D_f(g) = \emptyset$.) If we were to define

$$d_f^*(g_1,g_2) := |\inf D_f(g_2) - \inf D_f(g_1)|,$$

we would then find $d_f^*(g,\overline{S}f) = \varpi$ for such f and g. Hence, d_f^* is not a pseudometric (since the value of a pseudometric must be finite).

Hence, ϵ_0 is used to force d_f to take finite values. It may be thought of as the maximal uncertainty to be tolerated, the motivation being that we want "good" approximations, i.e.,

 ϵ -approximations for small values of ϵ .

On the other hand, if for any $f \in F$ and $g \in G$, there is an $\epsilon \geq 0$ such that $g \in S(f,\epsilon)$ (i.e., the "distance" between any solution and any point in G is finite), then d_f^* is always finite. Hence d_f^* is a pseudometric, and (2.5) holds for all $\epsilon \geq 0$, with d_f^* replacing d_f .

Remark 2.2. It would be more satisfying to be able to say that

(2.10)
$$S(f,\varepsilon) = B(\overline{S}f,d_f,\varepsilon)$$

in the conclusion of Theorem 2.1. However, we cannot do this in general. To see this, let d be a pseudometric on G, where G is not a singleton, and let $S: F \to G$ be any operator. Define a solution operator $S: F \times \mathbb{R}^+ \to 2^G$ by

$$S(f,\epsilon) := \begin{cases} \{g \in G : d(g,\hat{S}f) < \epsilon\} & \text{if } \epsilon > 0 \\ \{\hat{S}f\} & \text{if } \epsilon = 0 \end{cases}.$$

Suppose there exists $\overline{S}: F \to G$ and a family $\{d_f : f \in F\}$ of pseudometrics such that (2.10) holds.

We first note that $\overline{S} = \hat{S}$. Indeed, since $d_f(\overline{S}f, \overline{S}f) = 0$, we have

$$\overline{S}f \in B(\overline{S}f,d_f,0) = S(f,0) = {\hat{S}f},$$

i.e., $\overline{S}f = \hat{S}f$.

Next, we show that for any $f \in F$ and $g \in G$,

(2.11)
$$d_{f}(g,\overline{S}f) \leq d(g,\overline{S}f).$$

Indeed, let $\delta > 0$. Since $d(g,\overline{S}f) < d(g,\overline{S}f) + \delta$ and $Sf = \overline{S}f$, we have

$$g \in S(f,d(g,\overline{S}f) + \delta) = B(\overline{S}f,d_f,d(g,\overline{S}f) + \delta),$$

so that

$$d_f(g,\overline{S}f) \leq d(g,\overline{S}f) + \delta$$
.

Since $\delta > 0$ is arbitrary, (2.11) follows.

We claim that there exist $f \in F$ and $g \in G$ such that $d_f(g,\overline{S}f) > 0.$

Indeed, if $d_f(g,\overline{S}f)=0$ for all $f\in F$ and $g\in G$, we would have

$$g \in B(\overline{S}f, d_f, 0) = S(f, 0) = {\overline{S}f},$$

which would mean that

$$g = \overline{S}f \quad \forall f \in F, g \in G.$$

Fixing f and letting g vary, this would imply that G is a singleton, a contradiction.

Finally, we choose f \in F and g \in G such that (2.12) holds. Then $d_f(g,\overline{S}f) \ \underline{<} \ d_f(g,\overline{S}f)$ yields

$$g \in B(\overline{S}f, d_f, d_f(g, \overline{S}f)) = S(f, d_f(g, \overline{S}f)).$$

Since (2.12) holds, we have

$$d(g,\overline{S}f) < d_f(g,\overline{S}f),$$

contradicting (2.11).

3. A Family of Pseudometrics is Necessary

In this section, we reconsider Example 2.2, the solution of nonlinear equations. We show explicitly how to construct a family of pseudometrics such that (2.10) holds for all $\epsilon \geq 0$. Moreover, we show that a <u>family</u> of pseudometrics is necessary.

We first define $\overline{S}: F \to \mathbb{R}$ by letting $\overline{S}f$ be the zero of f that is smallest in magnitude; if there are two such zeros, choose the positive one. (Such a zero exists because f is continuous.)

Next, define $d_f : \mathbb{R} \times \mathbb{R} \to \mathbb{R}^+$ (for $f \in F$) by

$$d_{f}(x,y) := |f(x) - f(y)|$$
.

Then df is a pseudometric.

We then have

Theorem 3.1. $S(f,\epsilon) = B(\overline{S}f,d_f,\epsilon) \quad \forall f \in F, \epsilon \geq 0.$ Proof: Let $f \in F, \epsilon \geq 0$. Since $f(\overline{S}f) = 0$, we have

$$x \in S(f,\epsilon) \Leftrightarrow |f(x)| \leq \epsilon$$

$$\Leftrightarrow |f(x) - f(\overline{S}f)| \leq \epsilon$$

$$\Leftrightarrow d_f(x,\overline{S}f) \leq \epsilon$$

$$\Leftrightarrow x \in B(\overline{S}f,d_f,\epsilon).$$

Hence, Example 2.2 generates a family $\{d_{\mathbf{f}}: \mathbf{f} \in \mathbf{F}\}$ of pseudometrics, and the $d_{\mathbf{f}}$ -ball of radius ϵ about $\overline{\mathbf{S}}\mathbf{f}$ is the set of ϵ -approximations to the zeroset of \mathbf{f} , for any $\mathbf{f} \in \mathbf{F}$ and $\epsilon \geq 0$.

We now show that Example 2.2 cannot be generated by a single pseudometric.

Theorem 3.2. There does not exist an operator $\overline{S}:F\to\mathbb{R}$ and a single pseudometric d on \mathbb{R} such that

$$(3.1) S(f,\varepsilon) = B(\overline{S}f,d,\varepsilon) \forall f \in F, \varepsilon > 0.$$

<u>Proof</u>: Suppose there exists \overline{S} and d such that (3.1) holds. We first note that $\overline{S}f \in Z(f)$ for all $f \in F$. To see this, let $f \in F$. Since $d(\overline{S}f,\overline{S}f) = 0$, we have $\overline{S}f \in B(\overline{S}f,d,0) = S(f,0)$. Hence $|f(\overline{S}f)| \leq 0$ by (2.3), so that $f(\overline{S}f) = 0$ and $\overline{S}f \in Z(f)$, as claimed.

We next claim that

(3.2)
$$d(x,\overline{S}f) = |f(x)| \quad \forall x \in \mathbb{R}, f \in F.$$

Indeed, let $f \in F$, and $x \in \mathbb{R}$. Using (2.3) and (3.1), we have

$$|f(x)| \le |f(x)| \Rightarrow x \in S(f,|f(x)|) = B(\overline{S}f,d,|f(x)|)$$

 $\Rightarrow d(x,\overline{S}f) \le |f(x)|,$

and

$$d(x,\overline{S}f) \leq d(x,\overline{S}f) \Rightarrow x \in B(\overline{S}f,d,d(x,\overline{S}f)) = S(f,d(x,\overline{S}f))$$

$$\Rightarrow |f(x)| \leq d(x,\overline{S}f),$$

yielding (3.2).

We now let $x,y \in \mathbb{R}$ with $x \neq y$. Define $f_{\alpha} \in F$ by

(3.3)
$$f_{\alpha}(t) := \alpha(t - y) \quad \forall \alpha \in \mathbb{R}.$$

Then the first paragraph of this proof yields

$$\overline{S}f_{\alpha} \in Z(f_{\alpha}) = \{y\}$$
 $\forall \alpha \in \mathbb{R}$,

i.e.,

$$\overline{S}f_{\alpha} = y \qquad \forall \alpha \in \mathbb{R}.$$

Hence (3.4), (3.2), and (3.3) yield

$$d(x,y) = d(x,\overline{S}f_{\alpha})$$

$$= |f_{\alpha}(x)|$$

$$= |\alpha(x - y)|$$

$$= |\alpha||x - y|$$

Since $x \neq y$, this means that d(x,y) must be multiple-valued, a contradiction.

Note that the proof of Theorem 3.2 did not require the use of functions with multiple zeros. Hence, even if we consider Example 2.2 with F replaced by

 $F':= \{f: \mathbb{R} \to \mathbb{R} \mid f \text{ is continuous and has exactly one zero} \}$ we still cannot use a single pseudometric to generate this example.

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